# Fast solvers for implicit Runge-Kutta systems 

Jed Brown jedbrown@mcs.anl.gov<br>Debojyoti Ghosh ghosh@mcs.anl.gov

Mathematics and Computer Science Division
Argonne National Laboratory

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## Outline

The memory bandwidth problem

Implicit Runge-Kutta

Tensor product algebra

## Hardware Arithmetic Intensity

| Operation | Arithmetic Intensity (flops/B) |
| :--- | :---: |
| Sparse matrix-vector product | $1 / 6$ |
| Dense matrix-vector product | $1 / 4$ |
| Unassembled matrix-vector product | $\approx 8$ |
| High-order residual evaluation | $>5$ |


| Processor | Bandwidth (GB/s) | Peak (GF/s) | Balanced AI (F/B |
| :--- | ---: | ---: | ---: |
| E5-2680 8-core | 38 | 173 | 4. |
| Magny Cours 16-core | 49 | 281 | 5. |
| Blue Gene/Q node | 43 | 205 | 4. |
| Tesla M2090 | 120 | 665 | 5. |
| Kepler K20Xm | 160 | 1310 | 8. |
| Xeon Phi SE10P | 161 | 1060 | 6. |

## Optimizing Sparse Mat-Vec

■ Order unknowns so vector reuses cache (Cuthill-McKee)

- Optimal: $\frac{(2 \text { flops)(bandwidth) }}{\text { sizeof(Scalar)+sizeof(Int) }}$

■ Usually improves strength of ILU and SOR
■ Coalesce indices for adjacent rows (Inodes)

- Optimal: $\frac{(2 \text { flops)(bandwidth) }}{\text { sizeof (Scalar)+sizeof(Int)/i }}$
- Can do block SOR (much stronger than scalar SOR)

■ Default in PETSc, turn off with -mat_no_inode
■ Requires ordering unknowns so that fields are interlaced, this is (much) better for memory use anyway
■ Use explicit blocking, hold one index per block (BAIJ format)
■ Optimal: $\frac{(2 \text { flops)(bandwidth) }}{\text { sizeof(Scalar) }+ \text { sizeof (Int) } / b^{2}}$

- Block SOR and factorization

■ Symbolic factorization works with blocks (much cheaper)

- Very regular memory access, unrolled dense kernels

■ Faster insertion: MatSetValuesBlocked()

## This is a dead end

- Arithmetic intensity $<1 / 4$

■ Idea: multiple right hand sides

$$
\frac{(2 k \text { flops)(bandwidth) }}{\text { sizeof (Scalar) }+ \text { sizeof (Int) }}, \quad k \ll \text { avg. nz/row }
$$

■ Problem: popular algorithms have nested data dependencies
■ Time step
Nonlinear solve
Krylov solve
Preconditioner/sparse matrix
■ Cannot parallelize/vectorize these nested loops

Attempt: s-step methods in 3D


■ Amortizing message latency is most important for strong-scaling

- $s$-step methods have high overhead for small subdomains
- Limited choice of preconditioners (none optimal)

Attempt: space-time methods (multilevel SDC/Parareal)


■ PFASST algorithm (Emmett and Minion, 2013)
■ Zero-latency messages (cf. performance model of $s$-step)
■ Spectral Deferred Correction: iterative, converges to IRK (Gauss, Radau, ...)

- Stiff problems use imolicit basic intearator (svnchronizing on spatial


## Problems with SDC and time-parallel


c/o Matthew Emmett, parallel compared to sequential SDC
■ Number of iterations is not uniform, efficiency starts low

- Arithmetic intensity unchanged

■ Parabolic space-time (Greenwald and Brandt/Horton and Vandewalle)

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## Runge-Kutta methods

$$
\begin{gathered}
\dot{u}=F(u) \\
\underbrace{\left(\begin{array}{c}
y_{1} \\
\vdots \\
y_{s}
\end{array}\right)}_{Y}=u^{n}+h \underbrace{\left[\begin{array}{ccc}
a_{11} & \cdots & a_{1 s} \\
\vdots & \ddots & \vdots \\
a_{s 1} & \cdots & a_{s s}
\end{array}\right]}_{A} F\left(\begin{array}{c}
y_{1} \\
\vdots \\
y_{s}
\end{array}\right) \\
u^{n+1}=b^{T} Y
\end{gathered}
$$

- General framework for one-step methods

■ Diagonally implicit: $A$ lower triangular, stage order $\leq 2$
■ Singly diagonally implicit: all $A_{i j}$ equal, reuse solver setup, stage order $\leq 1$

- If $A$ is a general full matrix, all stages are coupled, "implicit RK"


## Implicit Runge-Kutta

| $\frac{1}{2}-\frac{\sqrt{15}}{10}$ | $\frac{5}{36}$ | $\frac{2}{9}-\frac{\sqrt{15}}{15}$ | $\frac{5}{36}-\frac{\sqrt{15}}{30}$ |
| :---: | :---: | :---: | :---: |
| $\frac{1}{2}$ | $\frac{5}{36}+\frac{\sqrt{15}}{24}$ | $\frac{2}{9}$ | $\frac{5}{36}-\frac{\sqrt{15}}{24}$ |
| $\frac{1}{2}-\frac{\sqrt{15}}{10}$ | $\frac{5}{36}+\frac{\sqrt{15}}{30}$ | $\frac{2}{9}+\frac{\sqrt{15}}{15}$ | $\frac{5}{36}$ |
|  | $\frac{5}{18}$ | $\frac{4}{9}$ | $\frac{5}{18}$ |

■ Implicit Runge-Kutta methods have excellent accuracy and stability properties

- Gauss methods with $s$ stages
- order $2 s,(s, s)$ Padé approximation to the exponential
- A-stable, symplectic

■ Radau (IIA) methods with $s$ stages
■ order $2 s-1, A$-stable, $L$-stable

- Lobatto (IIIC) methods with $s$ stages

■ order $2 s-2, A$-stable, $L$-stable, self-adjoint

- Stage order $s$ or $s+1$


## Method of Butcher (1976) and Bickart (1977)

■ Newton linearize Runge-Kutta system

$$
Y=u^{n}+h A F(Y)
$$

■ Solve linear system with tensor product operator

$$
S \otimes I_{n}+I_{s} \otimes J
$$

where $S=(h A)^{-1}$ is $s \times s$ dense, $J=-\partial F(u) / \partial u$ sparse
■ SDC (2000) is Gauss-Seidel with low-order corrector
■ Butcher/Bickart method: diagonalize $S=X \wedge X^{-1}$
■ $\Lambda \otimes I_{n}+I_{s} \otimes J$
■ $s$ decoupled solves
■ Problem: $X$ is exponentially ill-conditioned wrt. $s$

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## MatTAIJ: "sparse" tensor product matrices

$$
G=I_{n} \otimes S+J \otimes T
$$

- More general than multiple RHS (multivectors)
- Compare to multiple right hand sides in row-major
- Runge-Kutta systems have $T=I_{s}$ (permuted from Butcher method)

■ Stream $J$ through cache once, same efficiency as multiple RHS

128 nodes, 16 procs/node, small diffusion problem Method order nsteps time

| Gauss 4 | 8 | 10 | $3.4345 \mathrm{e}-01$ |
| :--- | :--- | :--- | :--- |
| Gauss 2 | 4 | 20 | $7.6320 \mathrm{e}-01$ |
| Gauss 1 | 2 | 40 | $1.1052 \mathrm{e}+00$ |

## Calibration and accuracy

■ Splitting errors plague multi-physics simulation
■ Verlet (leapfrog) integration is popular: symplectic and cheap
■ Stability problems: damping and even/odd decoupling
■ Models calibrated to compensate

- Force parametrizations in molecular dynamics
- Atmospheric column physics


## Impact of time step on autoconversion vs accretion partitioning (from Hui)

Global Mean Normalized w.r.t. Default Model Configuration



c/o Peter Caldwell (LLNL)
■ Models calibrated for "efficient" time step
■ Not longer solving the PDEs we write down
■ Many FTE-years to recalibrate when discretization changes

- Calibration eats up a big chunk of the IPCC policy timeline


## Implicit Runge-Kutta for advection

Table : Total number of iterations (communications or accesses of $J$ ) to solve linear advection to $t=1$ on a 1024-point grid using point-block Jacobi preconditioning of implicit Runge-Kutta matrix. The relative algebraic solver tolerance is $10^{-8}$.

| Family | Stages | Order | Iterations |
| :--- | ---: | ---: | ---: |
| Crank-Nicolson/Gauss | 1 | 2 | 3627 |
| Gauss | 2 | 4 | 2560 |
| Gauss | 4 | 8 | 1735 |
| Gauss | 8 | 16 | 1442 |

■ Naive centered-difference discretization

## Trade-offs in time integration

- Properties

■ Nonlinear stability (e.g., positivity preservation)
■ Stability along imaginary axis

- L-stability (damping at infinity)
- Implicitness and reuse

■ What is expensive?

- Function evaluation

■ Operator assembly/preconditioner setup
■ How much can be reused for how long?
■ Implicit solves
■ Can we find better solver algorithm?
■ More effort in setup?
■ What is "convergence"?

- Wave propagation: implicitness useless for convergence in a norm

■ Non-norm functionals could be robust

## Outlook

■ Next up: Algebraic multigrid for tensor product operators
■ Technicalities: imaginary rotation in coarse operators (cf. MG for Helmholtz)
■ Stochastic Galerkin have some structure
■ Is it possible to design methods with well-conditioned $S=X \wedge X^{-1}$

