

Topology-aware load balancing for parallel applications on multi-core systems and beyond

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Overview

Asymptotically optimal load balancing algorithm for multi-core machines.



Motivation

- Example of environment: CHARM++
 - Parallel task overdecomposition
 - Chares 1 2 3
 - Platform independent
 - Processing elements

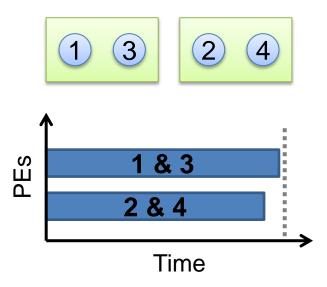


- Dynamic load balancing
 - Chare migration



Motivation: example

- Initial chare mapping
- Apply a load balancing algorithm
 - Based on data provided
 by CHARM++





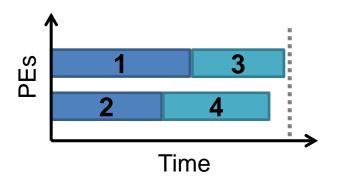


Motivation: example

- CHARM++ load balancing data
 - Communication graph
 - Current chare mapping
 - Chares' load





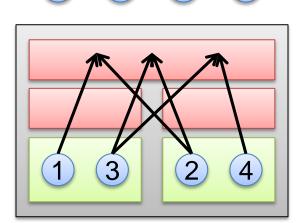


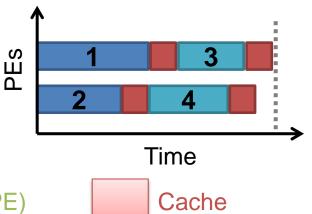




Motivation: example

- Missing information
 - Communication costs
 - Architectural information
 - Memory hierarchy



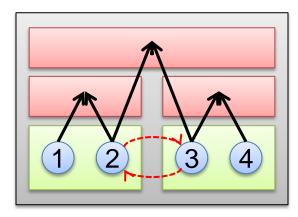


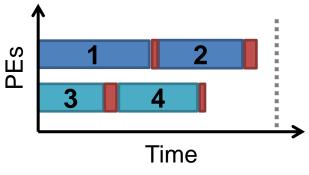
Processing Element (PE)



Motivation: example

- Knowledge about the memory hierarchy
 - Memory access costs
 - Reduce communication costs
 - Highly hierarchical systems
 - Even more on many-core





Cache



Processing Element (PE)



Motivation

- Our approach: **TopoLB**
 - Load balancing algorithm implemented on CHARM++
 - Combines *application information* and the *machine topology*
 - Works on UMA and NUMA machines
 - Is asymptotically optimal



Agenda

Motivation

ΤοροίΒ

Experiments And beyond



TOPOLB: idea

- Heuristic
 - Load balancing is NP-Hard
 - No initial assumption about the application
- Improve performance (reduce makespan)
 - By reducing unbalance
 - By reducing communication costs
 - While avoiding migrations (data movement costs)



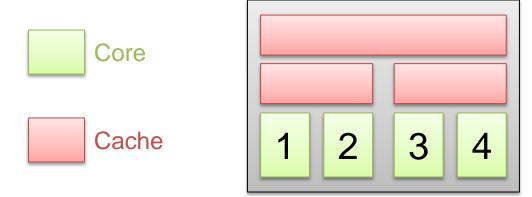
TOPOLB: data

- Application information provided by CHARM++
 - Chares' load
 - Reduce the makespan
 - Communication graph
 - Reduce the communication overhead
 - Bring together communicating chares
 - Current chare mapping
 - Avoid migration overheads



TOPOLB: data

- *Machine topology* our library
 - Memory hierarchy
 - Cache and memory sharing among cores
 - Memory access latencies
 - Estimate the communication time of each message through the memory hierarchy

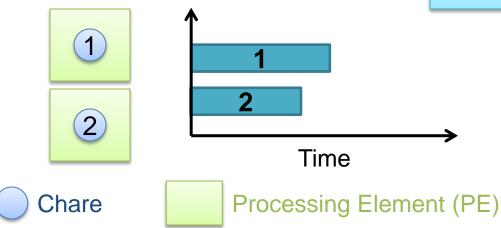




TOPOLB: algorithm

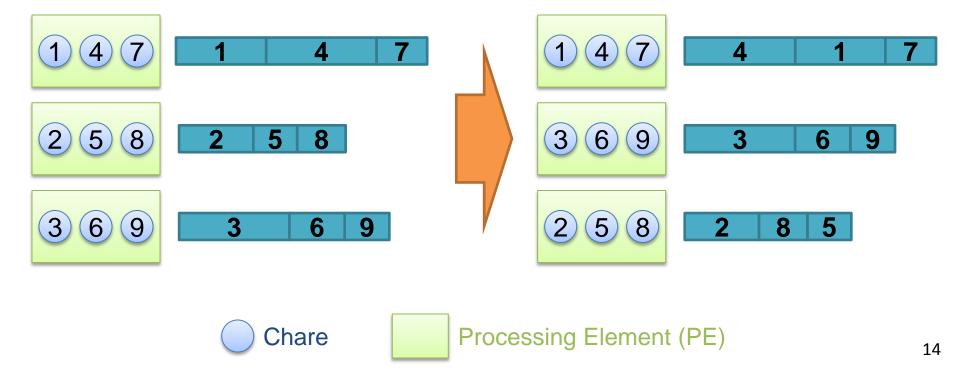
- Each chare has a load based on
 - Computation time
 - Communication time
 - Number of messages exchanged
 - Memory latencies

Communicating chares have a higher memory affinity



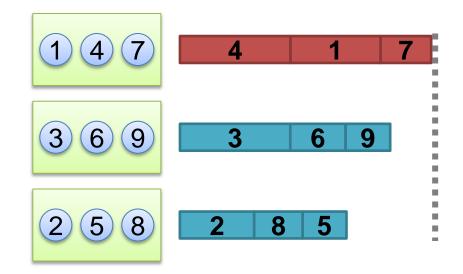


- Initial mapping
- Sort PEs (processors) based on their load



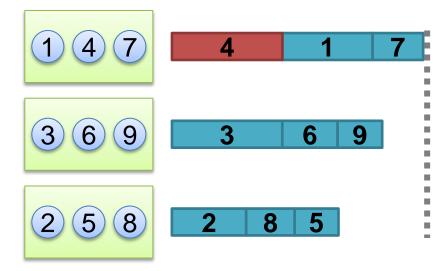


- Tries to migrate chares from the most loaded processor with a probability α
 - Or chooses another processor uniformly



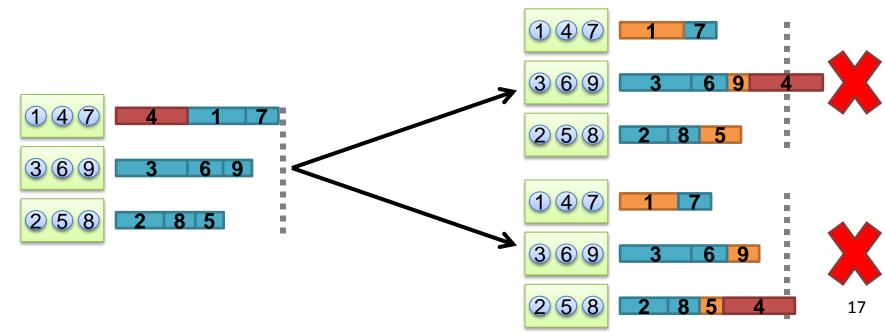


- Starts from the heaviest chare with a probability β
 - Or chooses another chare uniformly
 - From heaviest chare to lightest chare



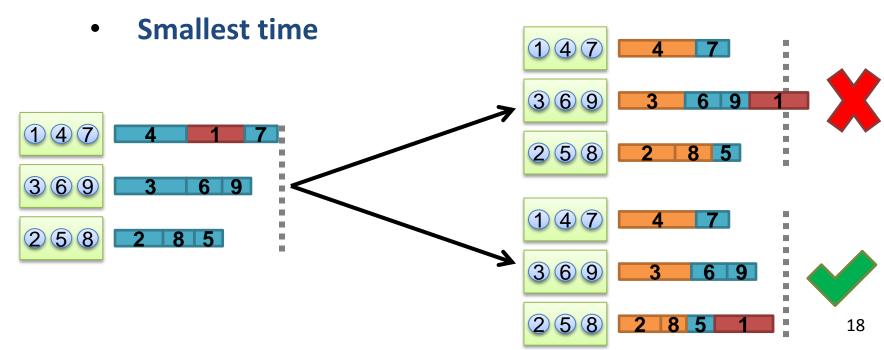


- A migration might decrease the makespan
- A migration may affect other chares
 - Communication time may change



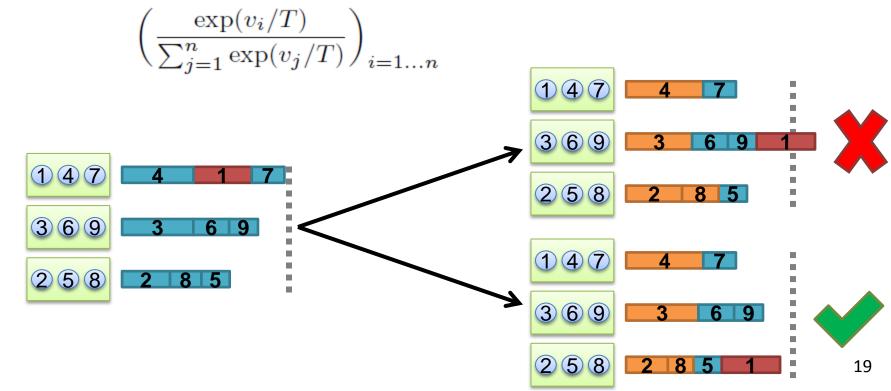


- Evaluates all possible processors
 - Migrates the chare to the processor that minimizes the makespan with high probability



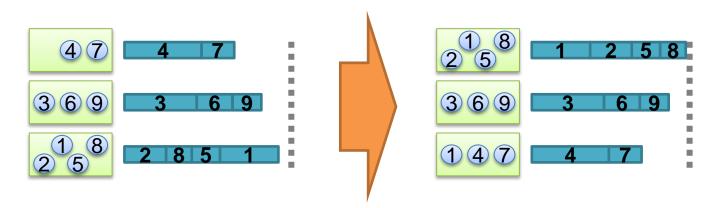


- High probability
 - Using a Gibbs distribution with a temperature T



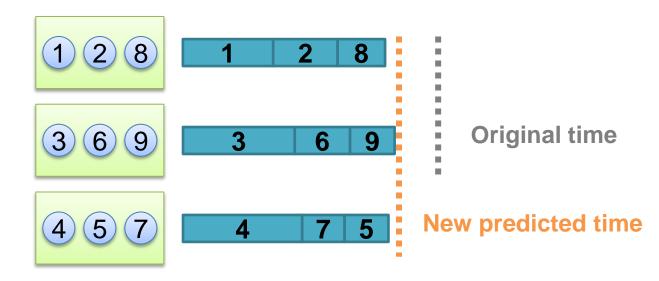


- Continues until no chare migrates from the heaviest processor
 - Cannot reduce the makespan
 - Bound to the heaviest processor





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 - Cannot reduce the makespan
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TOPOLB: implementation details

• Exponential backoff

- Reduce the load balancing overhead by not computing the algorithm on all LB calls
- Launched when **no migrations happen**





Agenda

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TOPOLB

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And beyond

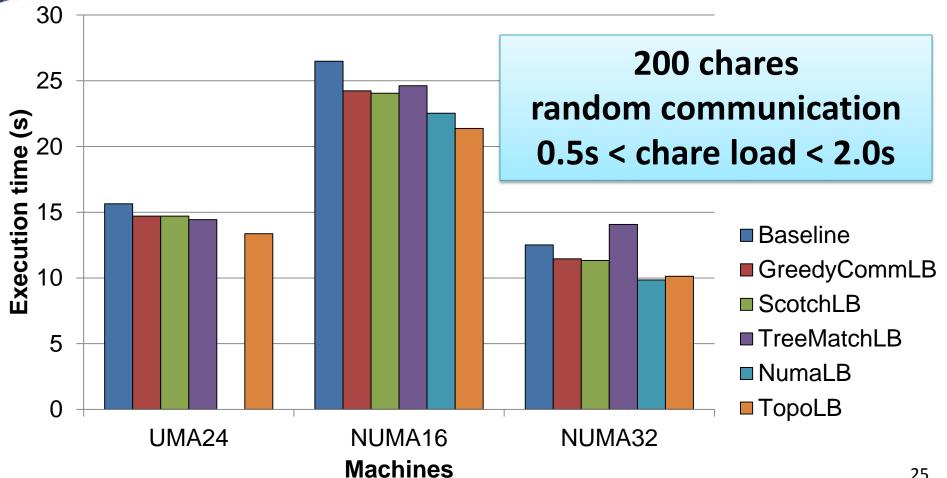


Experiments

- 3 different machines: UMA24, NUMA16 and NUMA32
- 4 other load balancers: GREEDYCOMMLB, SCOTCHLB, TREEMATCHLB and NUMALB
- 2 benchmarks: lb_test and mol3D
- α , β close to 1, T close to 0

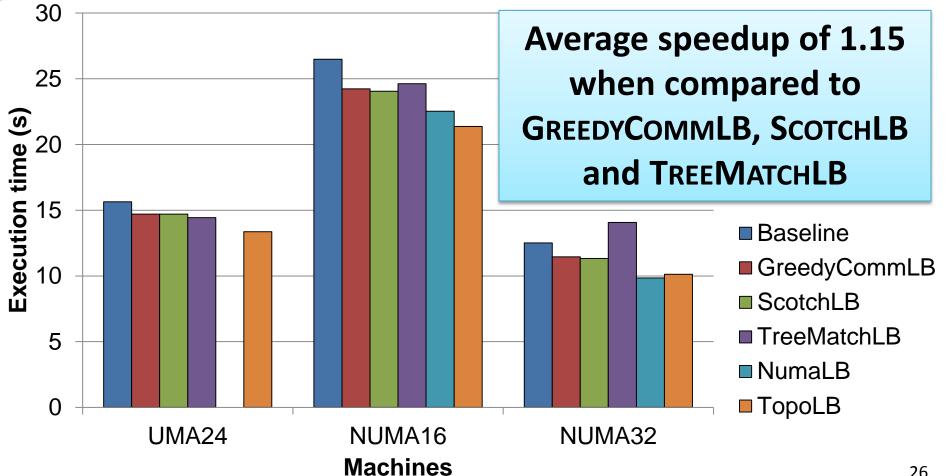


lb_test



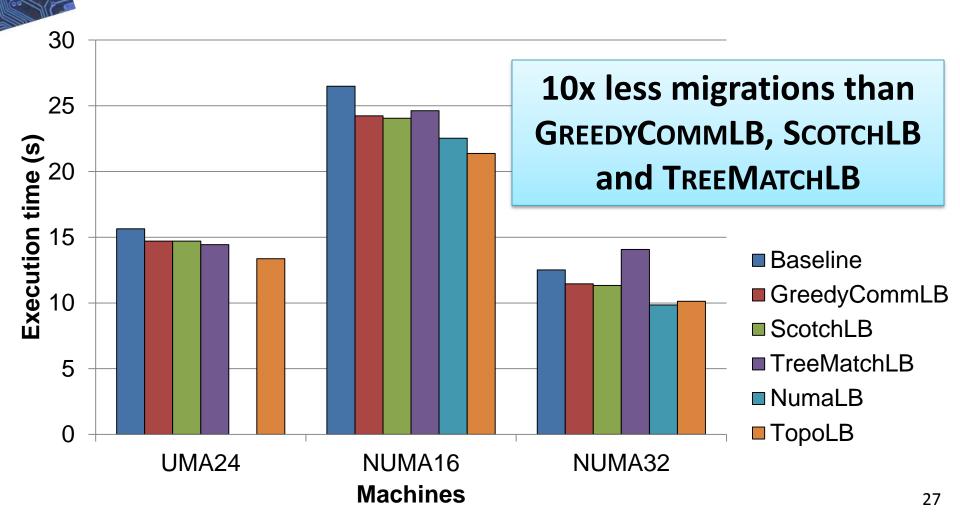


lb_test





lb_test





mol3D

- Apolipoprotein-A1
- No load balancer improved performance
- **TOPOLB** performed **30x less migrations** than other load balancers
 - But took 2 to 3x more time on its load balancing decisions



mol3D

- View using Projections
 - 3x20 iterations, 3 calls to TOPOLB
 - ~7000 chares





mol3D

Application starts balanced





mol3D

TOPOLB improves the iteration time





mol3D

Iteration time diverges strongly





Agenda

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TOPOLB

Experiments

And beyond



And beyond (concluding remarks)

- TOPOLB presented a
 - Small number of migrations
 - Fast convergence
 - High computational cost (overhead)
 - Reduced by the exponential backoff



And beyond

- Working to extend TOPOLB to clusters of multi-core machines
 - Two levels: Another LB for the cluster, TOPOLB for the compute nodes
 - TOPOLB can be **too costly** for large machines and applications



And beyond

- Working to **better understand and evaluate** the behavior of CHARM++ **load balancers**
 - Debugging library
 - Use real applications

• Added part of our machine model to HWLOC



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Thank you.

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